# **Artificial Intelligence**

## Lecture 21

## Prolog Programming for AI



#### **Prepared by:**

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**Prolog Programming for AI** 

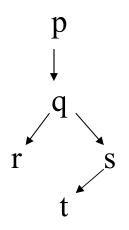
# **Lecture Outlines**

- Prolog Programs:
  - Prolog Programs on Tree and Graph Problems.
  - Prolog Programs for DFA and NFA



## **Example-20:** Prolog program for Tree...

• Consider acyclic directed graph:



• Graph G is represented by a set of connected edges:

**G** = {edge(p, q), edge(q, r), edge(q, s), edge(s, t)}

• Write Prolog program to check whether there is a route from one node to another node.

### **Example-20:** Prolog program for Tree

- Define Route:
- Route from node A to B is defined as follows:
  - Route from A to B, if there is an edge from A some node C and a route from C to B.
  - Route from A to B, if there is a direct edge from A to B.

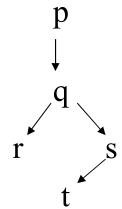
#### • Prolog Program:

- edge(p, q).
- edge(q, r).
- edge(q, s).

edge(s, t).

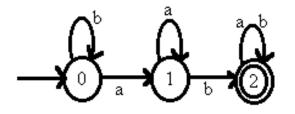
route(A, B) :- edge (A, C), route(C, B).

route(A, B) :- edge(A, B).



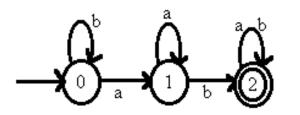
• Query: Check whether there exist a route between p and t.

- When a state table program are loaded into Prolog, the parser is used to check whether inputs to the DFA are acceptable or not.
- Consider a state diagram for a DFA that accepts the language (a,b)\*ab(a,b)\* is as follows:

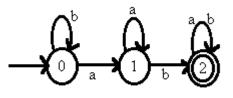


• Write a program to simulate a parser for an arbitrary deterministic finite automaton (DFA).

- In Prolog, an automaton can be specified by three relations:
  - 1) A unary relation 'start' which defines the initial state of the automation;
  - 2) A unary relation '**final**' which defines the final states of the automation;
  - A three-argument relation 'delta' which defines the state transitions so that: delta(S1, X, S2); this means that a transition from S1 to S2 is possible when the current input symbol X is read.



- Prolog code:
- Facts:
  - start(0).
  - delta(0,a,1).
  - delta(0,b,0).
  - delta(1,a,1).
  - delta(1,b,2).
  - delta(2,a,2).
  - delta(2,b,2).
  - final(2).



- The simulator is programmed as a unary relation parse (L) and a binary relation trans(S,L):
- The parse (L) relation can be defined by:
  - 1) Starting from the initial state S and calling the relation trans(S,L).
- The trans (S, L) relation can be defined by two clauses:
  - 1) The empty string, [], is accepted from a state S if S is a final state.
  - 2) A non-empty string is accepted from S if reading the first symbol in the string can bring the automation into some state, S1 and the rest of the string is accepted from S1.

#### • **Prolog Rules:**

- Suppose that both the driver program and the state table program are loaded ...
- ?- parse([b,b,a,a,b,a,b]).
  0 [b,b,a,a,b,a,b]
  0 [b,a,a,b,a,b]
  0 [a,a,b,a,b]
  1 [a,b,a,b]
  1 [b,a,b]
  2 [a,b]
  2 [b]
  0 [b,b,a]
  0 [b,a]
  - 2 [] 0 [a] yes no

- Simulating a NFA in Prolog, an automaton can be specified by three relations:
  - 1) A unary relation '**final**' which defines the final states of the automation;
  - 2) A three-argument relation '**trans**' which defines the state transitions so that: **trans(S1, X, S2)**

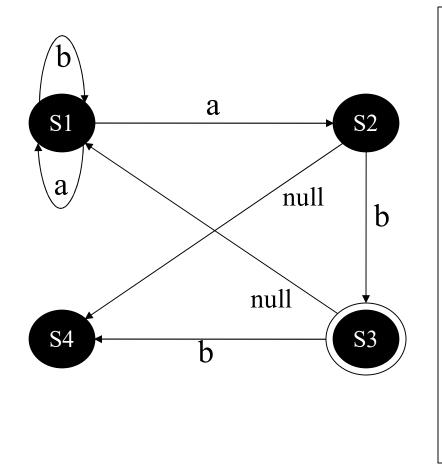
This means that a transition from S1 to S2 is possible when the current input symbol X is read;

3) A binary relation 'silent'

#### Silent(S1, S2)

Meaning that a silent move is possible from state S1 to S2.

• Consider the following non-deterministic finite automaton:



- For the given automation, the three relations are:
  - final(S3).
  - trans(S1, a, S1).
  - trans(S1, a, S2).
  - trans(S1, b, S1).
  - trans(S2, b, S3).
  - trans(S3, b, S4).
  - silent(S2, S4).
  - silent(S3, S1).

- We will represent input strings as Prolog lists. So, the string 'aab' will be represented by [a, a, b].
- The simulator is programmed as a binary relation, **accepts(S, X)**.
- The 'accepts' relation can be defined by three clauses:
  - 1) The empty string, [], is accepted from a state **S** if State is a final state.
  - 2) A non-empty string X is accepted from S if reading the first symbol in the string X can bring the automation into some state, S1 and the rest of the string is accepted from S1.
  - 3) A string is accepted from State if the automation can make a silent move from S to S1 and then accept the whole input string from S1.

• These rules can be translated into Prolog as:

```
accepts(String, []) :- % Accept empty string
    final(S).
accepts(S, [X | R]) :- % Reading 1<sup>st</sup> symbol
    trans(S, X, S1), accepts(S1, R).
accepts(S, String) :- % Silent move
    silent(S, S1), accepts(S1, String).
```

- Query:
- Acceptance of the input string *aaab*:

?- accepts(s1, [a,a,a,b]).

yes

- Initial state for input string *ab*:
  - ?- accepts(S, [a,b]).

$$\mathbf{S}=\mathbf{s1};$$

$$S = s3$$

• Find the string of length 3 that are accepted from state1.

X1 = a	X2 = a	X3 = b;
X1 = b	X2 = a	X3 = b;

**Prolog Programming for AI** 

# THE END